

Bounded lookahead in quantity insensitive stress assignment

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Introduction

- In iterative stress languages, stress is placed on every second or third syllable in the word
 - $\sigma\sigma\sigma\sigma\sigma\sigma\sigma\to \acute{\sigma}\sigma\acute{\sigma}\sigma\acute{\sigma}\sigma$
- What is the best characterization of these kinds of patterns?
- What formal properties do they share?

Introduction

- Formal language theory (FLT) delineates classes of functions that serve as typological hypotheses for stress assignment
- Some previous "big-picture" FLT work on stress as a function (Hao & Andersson 2019; Koser & Jardine 2020)
- No treatment of iterative stress patterns in particular

Results

- Typological split among iterative patterns less complex *output strictly* local (OSL) patterns¹ and more complex subsequential patterns²
- All more complex patterns share property of "look-ahead" despite surface differences
- If you separate the iteration of stress from the look-ahead, iterative patterns look the same
- More restrictive characterization of iterative stress

¹ Chandlee & Heinz (2018) ² Mohri (1997)

Why this matters

- A step closer to answering the question what is the proper characterization of stress as a function $?^1$
- Unites a group of patterns with surface differences based on their computational properties
- Bears on extra metricality 2 and non-finality 3
- Informative with regard to decomposition of subsequential functions

 1 Koser & Jardine (2020) 2 Liberman & Prince (1977) 3 Prince & Smolensky (1993)

Plan

- Background (FLT, stress)
- Complexity of the patterns
- Address the more complex patterns
- Implications

Complexity



- FLT complexity classes divide space of possible functions based on expressive power of those functions
- Phonology is *regular* (Johnson 1972; Kaplan & Kay 1994)
- In fact, most is subregular (Rogers et al. 2013; Heinz 2018)

FLT and phonology

- Classes correspond to different phonological patterns; different information
- Input strictly local (ISL) functions 1 bounded information in the input
 - Ex: initial stress: $\#\sigma\sigma\sigma\sigma\sigma \to \#\sigma\sigma\sigma\sigma$
- Output strictly local (OSL) functions² bounded information in the output
 - Ex: binary stress: $\#\sigma\sigma\sigma\sigma\sigma\sigma \rightarrow \#\sigma\sigma\sigma\sigma\sigma$
- Subsequential functions ^3 - information up to current input symbol, more powerful
 - Ex: LHOR stress: $LLHL \rightarrow LLHL \rightarrow LLLLL \rightarrow LLLLL$

 $^{^1}$ Chandlee (2014) 2 Chandlee & Heinz (2018) 3 Mohri (1997)

Stress

- Iterative binary quantity insensitive (QI) stress
 - binary: Murinbata¹: $\delta\sigma, \delta\sigma\delta, \delta\sigma\delta\sigma, \delta\sigma\delta\sigma\delta, \delta\sigma\delta\sigma\delta\sigma, \delta\sigma\delta\sigma\delta\sigma\delta$...
 - non-finality: Pintupi²: $\delta\sigma, \delta\sigma\sigma, \delta\sigma\delta\sigma, \delta\sigma\delta\sigma\sigma, \delta\sigma\delta\sigma\delta\sigma, \delta\sigma\delta\sigma\delta\sigma$...
- What is the complexity of different iterative stress?

 1 Street & Mollinjin (1981) $^{-2}$ Hansen & Hansen (1969) $^{-3}$ Kaye (1973) $^{-4}$ Furby (1974)

Stress

- Given iterative patterns are OSL, adopt it as null hypothesis
- Stress as string-to-string mapping from input to output with finite state transducers (FST)
- Function classes have well-understood FST properties

Stress

• Properties of FSTs make properties of the function apparent

start
$$\longrightarrow q_0 \xrightarrow{\sigma: \hat{\sigma}} q_1 \longrightarrow \sigma: \sigma$$

initial stress:

$\sigma\sigma$	\rightarrow	$\sigma\sigma$
$\sigma\sigma\sigma$	\rightarrow	$\sigma \sigma \sigma$
$\sigma\sigma\sigma\sigma\sigma$	\rightarrow	<u>ό</u> σσσ
σσσσσ	\rightarrow	<u>ό</u> σσσσ
•••	\rightarrow	•••

Binary: OSL

- Murinbata: $\dot{\sigma}\sigma, \dot{\sigma}\sigma\dot{\sigma}\sigma, \dot{\sigma}\sigma\dot{\sigma}\sigma\dot{\sigma}\sigma, \dot{\sigma}\sigma\dot{\sigma}\sigma\dot{\sigma}\sigma, \dot{\sigma}\sigma\dot{\sigma}\sigma\dot{\sigma}\sigma\dot{\sigma}\sigma$...
- Placement of stress determined based on output, OSL



Non-finality: non-OSL

- Pintupi: *ό*σ, *ό*σσ, *ό*σ*ό*σ, *ό*σ*ό*σσ, *ό*σ*ό*σ*ό*σ, *ό*σ*ό*σ*ό*σ...
- Binary stress would stress final
- \bullet Every odd syllable, needs to know am I at the end of the word?
- Requires lookahead (seen as "waiting" $\lambda)$
- Lookahead not OSL, is subsequential

start
$$\rightarrow q_0 \xrightarrow{\sigma: \hat{\sigma}} q_1 \xrightarrow{\sigma: \sigma} q_2 \xrightarrow{\sigma: \lambda} q_3: \sigma$$

Clash: non-OSL

- Ojibwe: $\sigma \dot{\sigma}, \sigma \dot{\sigma} \dot{\sigma}, \sigma \dot{\sigma} \sigma \dot{\sigma}, \sigma \dot{\sigma} \sigma \dot{\sigma} \dot{\sigma}, \sigma \dot{\sigma} \sigma \dot{\sigma} \sigma \dot{\sigma}, \sigma \dot{\sigma} \sigma \dot{\sigma} \dot{\sigma} \dot{\sigma}$...
- Binary stress would miss final stress
- Not OSL, is subsequential



Internal lapse: non-OSL

- Binary stress (right to left) would stress penult
- Not OSL, is subsequential

Taking stock

- Non-fin, clash, internal lapse patterns all subsequential
- All share property of lookahead
- Despite surface differences, similar computational properties

Two functions

- Capture the similarity by separating the iteration of stress from the lookahead
- One OSL function that blindly iterates binary stress
- One ISL function that acts like lookahead by "cleaning up"
- Output of OSL is input of ISL, like rule ordering

Two functions: Non-fin

 $\sigma\sigma\sigma\sigma\sigma\sigma \rightarrow \sigma\sigma\sigma\sigma\sigma\sigma \rightarrow \sigma\sigma\sigma\sigma\sigma\sigma$

- $\sigma\sigma\sigma\sigma\sigma\sigma\sigma \rightarrow \sigma\sigma\sigma\sigma\sigma\sigma\sigma \rightarrow \sigma\sigma\sigma\sigma\sigma\sigma$
- OSL function stresses every odd syllable left-to-right
- ISL function removes final stress if present
- Iteration is like Murinbata



Two functions: Clash

- $\sigma\sigma\sigma\sigma\sigma\sigma \rightarrow \sigma\sigma\sigma\sigma\sigma\sigma \rightarrow \sigma\sigma\sigma\sigma\sigma\sigma$
- Every even syllable
- Add final stress if not present

start
$$q_0 \xrightarrow{\sigma:\sigma} q_1$$
 start $q_0 \xrightarrow{\sigma:\sigma} q_1 \xrightarrow{\sigma:\sigma} \sigma:\sigma$

Two functions: internal lapse

- Every even syllable (right-to-left)
- Add initial stress if not present and delete peninitial stress if present

start
$$q_0 \xrightarrow{\sigma:\sigma} q_1$$
 start $q_0 \xrightarrow{\sigma:\sigma} q_1 \xrightarrow{\sigma:\sigma} q_2 \xrightarrow{\sigma:\sigma} \sigma:\sigma$
 $\sigma:\sigma \xrightarrow{\sigma:\sigma} \sigma:\sigma$

Summary

- Separating the lookahead from the iteration highlights underlying similarity in patterns with surface differences
- Creates mini typology: some delete, some add, some delete and add

Further Issues

- ISL function is not just any arbitrary ISL function; only ever needs one or two input symbols (similar for OSL function)
- State some restriction on them. Restriction on the interaction ala McCollum et al. (2018)?
- Restrictions very important otherwise difference with subsequential function is unclear
- Without restrictions, can *any* subsequential function be broken down in this way?

Further Issues

- Similar in spirit to extra metricality 1 and non-finality 2 analyses
- More like non-finality, all syllables remain in computation
- Obviates need for function reattaching extrametrical syllables trade off
- Neither of the above apply to clash or internal lapse cases

 $^{^1}$ Liberman & Prince (1977) $^{\ 2}$ Prince & Smolensky (1993)

Further Issues

- Can this be extended somehow to othe stress cases?
 - bidirectional: Cahuilla¹: Stress every other syllable in both directions, starting at the root-initial syllable
 - ternary?
 - quantity sensitive languages?

 $^{^{1}}$ Seiler (1977)

Thanks

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26